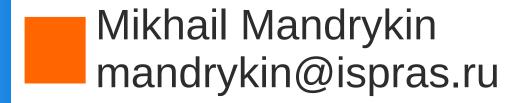
The ISoLA 2012 symposium Heraklion, Crete 15 Oct 2012



Pointer analysis with uninterpreted functions





Institute for System Programming of the Russian Academy of Sciences



Linux Driver Verification

Currently a model can be accurate if it *doesn't* make heavy use of (or rely on):

- Bitwise operations (e.g. &, |, ^)
- Arithmetic over/underflow error detection
- Non-linear arithmetic
- Multithreading (race conditions detection)
- Recursive data types (e.g. lists)
- Arrays
- Nested structures (and container_of macro)
- Pointer aliasing



Linux Driver Verification (2)

Bugfix example

```
(/drivers/connector/connector.c, commit 663dd6d, Linux 2.6.40 i.e. also 3.0):
        static int cn_call_callback(struct sk_buff *skb)
         {
             int err = -ENODEV;
             if (cbq != NULL) {
                 err = 0;
                 kfree_skb(skb);
                                                      Double freeing
             return err;
                                                    of a socket buffer
        }
                                                            (skb)
             err = cn_call_callback(skb);
             if (err < 0)
                 kfree_skb(skb);
```



Linux Driver Verification (3)

Typical example of rule with *inaccurate* model:

- Don't call kfree_skb twice
 - Avoid double freeing of skb buffers

- Socket buffers (or skbs) are represented by sk_buff structures
- Pointers to the structures are frequently located in arrays, lists, queues and other data structures
- This leads to many false positive (spurious UNSAFE) verdicts



Pointer analysis

BLAST with Andersen's alias analysis

In many cases it's really useful!

```
* kfree_skb(info->rx_skb);
info->rx_skb = NULL;
....
The infos here are may-aliased
    dtl1_receive(info);
    dtl1_receive(dtl1_info_t *info)
    {
        The call is
unreachable
        if (info->rx_skb)
```

kfree_skb(info->rx_skb);



Pointer analysis (2)

But still...

```
for (\underline{i} = 0; \underline{i} < RIONET_RX_RING_SIZE; \underline{i}++)
    kfree_skb(rnet->rx_skb[i]);
                                   Arrays
   RIONET_RX_RING_SIZE defaults to 128
while (db->rx_avail_cnt) {
    kfree_skb(db->rx_ready_ptr->rx_skb_ptr);
    db->rx_ready_ptr = db->rx_ready_ptr<u>->next</u>_rx_desc;
    db->rx avail cnt--;
                                                  Linked lists
   db->rx_avail_cnt <= RX_DESC_CNT == 32
nf_conntrack_put_reasm(skb->nfct_reasm);
```

Field to variable aliasing



Pointer analysis (3)

But still...

struct bdx_priv has 26 fields. 11 are other structures or pointers to structures, some of which have more than 25 fields, some of which are also structures... and we are to update every subfield in each of priv may-aliases when passing the parameter



Pointer analysis (4)

BLAST's and alternative approaches

Tool/approach	Pointers	Structures	Arrays	Recursive data structures	Pointer arithmetic	Performance
BLAST with "closure depth"	+	# (finite depth)	-	-	-	+
Optimized BLAST with "infinite closure depth"	+	+	-	-	-	+
BLAST with lazy shape analysis ("BLAST 3.0")	+	+	-	+	-	?
Bounded Model Checking	+	+	+	t(finite depth)	+	-
CPAchecker with predicate analysis(current implementation)	+	-	-	ı	-	+
Our approach (uninterpreted functions)	+	+	+	t (finite depth)	+	?



Suggested approach

What is the *initial* idea behind the approach?

Very simple, terribly inefficient but precise memory model:

Suggested approach (2)

Even more inefficient...

- Most state-of-the-art SMT-solvers still don't fully support array interpolation
- So we'll use uninterpreted functions
- **No** store (\cdot,\cdot,\cdot) operation
- Congruence: $a = b \rightarrow f(a) = f(b)$
- $m_1(a_1) = 1$, $m_2(a_2) = 2$, $m_2(a_1) = ?$
- We need to explicitly encode retention of earlier assigned values
- $\mathbf{a}_2 \neq \mathbf{a}_1 \rightarrow \mathbf{m}_2(\mathbf{a}_1) = \mathbf{m}_1(\mathbf{a}_1)$ and so on for every \mathbf{a}_i

Suggested approach (3)

How do we encode memory regions?

- One uninterpreted constant for each region
- Each region has positive address $(b_i>0)$
- Regions don't intersect:

$$B(b_i + k) = i$$
, $0 \le k < s$, where s is the size of the region

- $b_i + k = b_i + l \rightarrow B(b_i + k) = B(b_i + l) \rightarrow i = j$
- So number of such equalities is linear



Suggested approach (4)

What are suggested optimizations?

• *Typing* i.e. one array per one simple data type

Pure variables, i.e. variables that don't have aliases

```
int i; // just a counter
```

e.q. char *,

long int,

 Structure field assignment optimization, i.e. omitting the antecedents if offsets are known to be unequal in advance

// `&i' occurs **nowhere** in the code

```
e.g. updating Skb1->next can't influence any Skb2->prev though they have the same type
```



Suggested approach (5)

Further optimizatinos

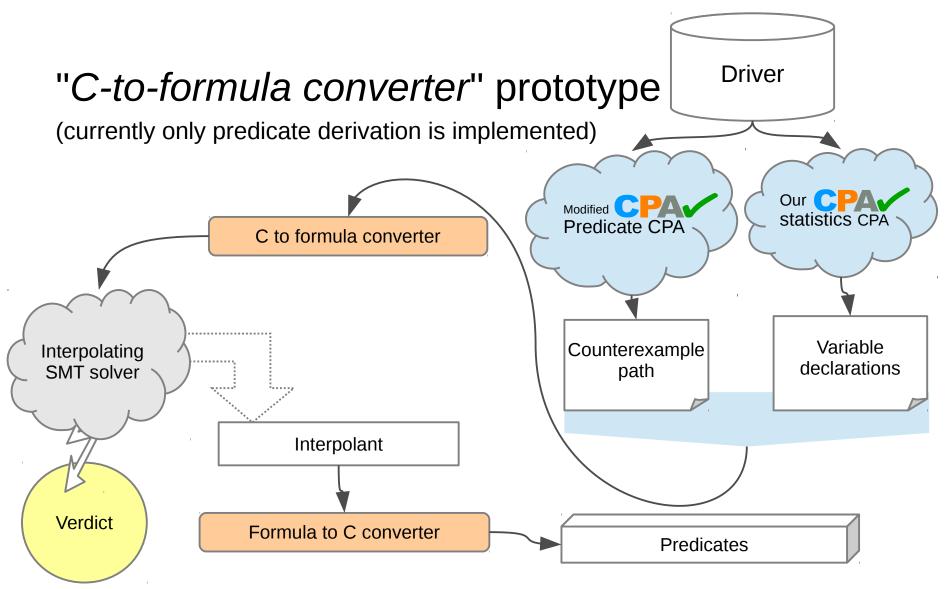
- Using constant subexpressions for initialization
 e.g. kzalloc(sizeof(*info), GFP_KERNEL)
- Amortization of sequential assignments

```
e.g. for (i = 0; i < MAX_SKB_FRAGS + 1; i++) {
    lwords = 7 + (i * 3);
    ... /* pad it with 1 lword */
    txd_sizes[i].qwords = lwords >> 1;
    txd_sizes[i].bytes = lwords << 2;
}
// No reading through any pointer during the entire loop
// So let's update the memory just once after the loop!</pre>
```

Applying preliminary alias analysis (again!)



Evaluation





Evaluation (1)

Pointer target statistics

		By type					
Driver	Total	1	2	3	4	5	By offset (max)
bluetooth/bpa10x.ko (2 skb)	422	unsigned char 134	signed char 112	unsigned long int 44	unsigned short int 26	signed long int 18	18
bluetooth/dtl1_cs.ko (32 skb)	3814	signed char 1609	unsigned char 639	unsigned long int 536	unsigned short int 312	signed long int 170	137
isdn/hysdn/hysdn.ko (20 skb)	2120	signed char 960	unsigned char 340	unsigned long int 300	unsigned short int 180	signed long int 100	80
hid/usbhid/usbkbd.ko (no rule model applied)	789	unsigned char 293	signed char 329	unsigned long int 39	struct list_head* 12	unsigned short int 12	11
net/usb/cdc-phonet.ko (no rule model applied)	224	unsigned long int 38	signed char 59	signed char * 16	unsigned char 16	struct list_head* 12	17



Evaluation (2)

Formula sizes in KB

Driver	No optimizations	Target filtering	Pure variables	Both
bluetooth/bpa10x.ko (2 skb)	448	369	311	220
bluetooth/dtl1_cs.ko (32 skb)	3700	2400	623	456
isdn/hysdn/hysdn.ko (20 skb)	726	474	101	95
hid/usbhid/usbkbd.ko (no rule model applied)	352	279	119	88
net/usb/cdc-phonet.ko (no rule model applied)	255	166	15	15



Evaluation (3)

Sample interpolant and predicates

```
true
(and
  (= usbpn_open!!i~1 0.0)
  (= usbpn_open!!dev~1 usbpn_open!!pnd~1))
(and
  (= usbpn_open!!i~1 0.0)
  (= usbpn_open!!dev~1 usbpn_open!!pnd~1))
(= (struct-urb-*~2 (+ (+ usbpn_close!!pnd~1 usbpn_close!!i~1) 66.0)) 0.0)
false
                                                      ~70 locations
                                                      ~60 operators in path
                                                      MathSAT interpolation time: 0.039s
0 == 0
usbpn_open::i == 0 && usbpn_open::dev == usbpn_open::pnd
usbpn_open::i == 0 && usbpn_open::dev == usbpn_open::pnd
usbpn_close::pnd->urbs[usbpn_close::i] == 0
0 < 0
```



Conclusions

- Approach isn't carefully evaluated yet
- Current results are not disappointing
- Not all optimizations are implemented → better performance expected
- Many low-level C language features are supported (such as pointer arithmetics, container_of macro etc.)
- We are planning further investigation of the approach and its implementation as a CPA in the CPAchecker(CPA✓) tool



Thank you!

Mikhail Mandrykin mandrykin@ispras.ru



Institute for System Programming of the Russian Academy of Sciences